

Reduction of Transcendental Decision Problems over the Reals

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July 18, 2024



Authors



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Shia”)

I am Rizeng Chen, the people on the left and I am currently a third-year PhD candidate under the supervision of Prof. Xia.

1 Introduction

2 History

3 Reduction

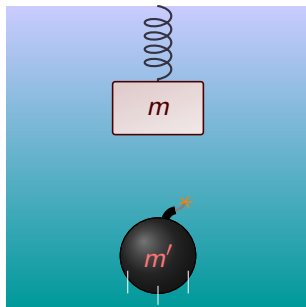
4 Proof Sketch

5 Multivariate Case

6 Implementation

Problems with Transcendental Constraints

A **mass** is attached to a **spring** in the water. A **bomb** is thrown beneath the mass. The bomb will **explode** when it hits the mass!



Question: Will the mass and the bomb collide at some time $t > 0$?

Problems with Transcendental Constraints

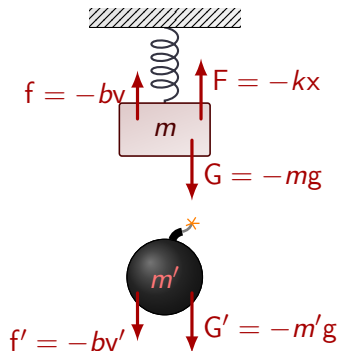


Figure 1: Free Body Diagram

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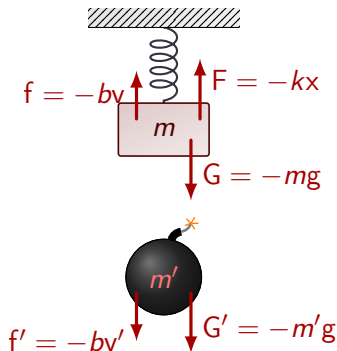


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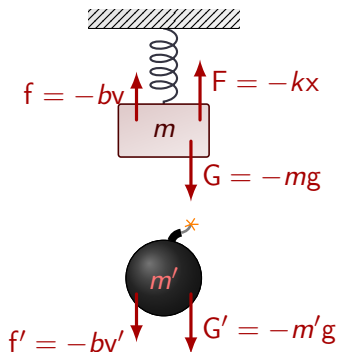


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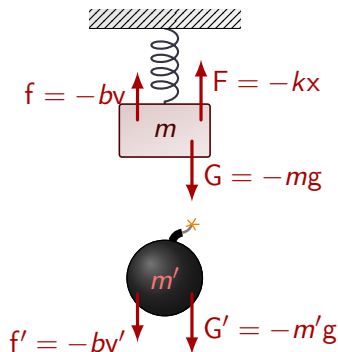


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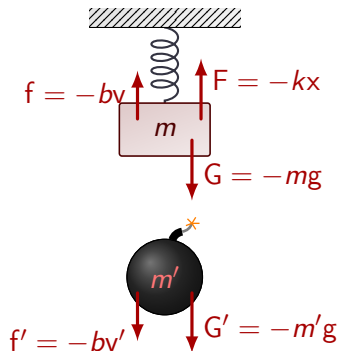


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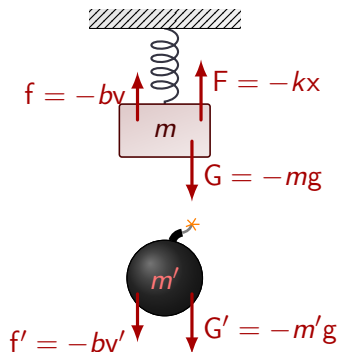


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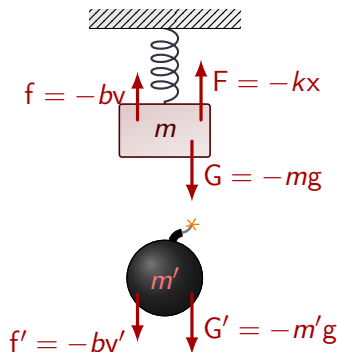


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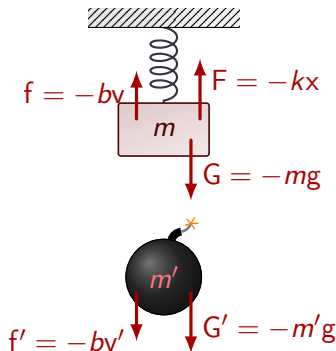


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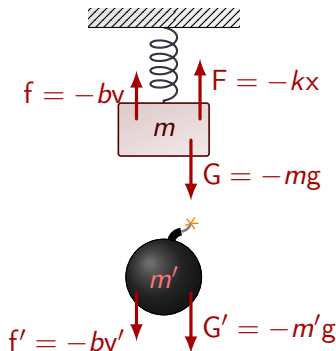


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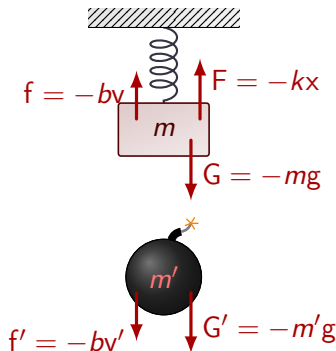


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- Suppose $m = m' = 1\text{kg}$, $k = 10\text{N/m}$, $b = 2\text{N} \cdot \text{s/m}$ and $g = 10\text{m/s}^2$.
- The initial positions are $x(0) = 0\text{m}$, $y(0) = -5\text{m}$, and the initial velocities are $\dot{x}(0) = -12\text{m/s}$ and $\dot{y}(0) = 9\text{m/s}$.

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$$\begin{cases} \ddot{x} &= -2\dot{x} - 10x - 10 \\ \dot{x}(0) &= -12 \\ x(0) &= 0 \end{cases}, \begin{cases} \ddot{y} &= -2\dot{y} - 10 \\ \dot{y}(0) &= 9 \\ y(0) &= -5 \end{cases}.$$

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- **Solving the ODE:**

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- **Transcendental constraints** naturally arise in the real world.
- How do we handle them?

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Polynomials are easy...

- **Tarski (1951)** showed that the first-order theory over a real closed field is decidable (polynomial equations and inequalities).

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- Schanuel's Conjecture (SC)

Suppose $z_1, \dots, z_n \in \mathbb{C}$ are linear independent over \mathbb{Q} , then

$$\text{tr.deg } \mathbb{Q}(z_1, \dots, z_n, e^{z_1}, \dots, e^{z_n}) \geq n$$

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- Achatz, McCallum and Weispfenning (2008) presented an algorithm to decide polynomial-exponential problems (later generalized to $\ln x$ and $\arctan x$).
- At the same time, Strzeboński (2008) studied the real root isolation of exp-log functions (then extended to tame elementary functions and exp-log-arctan functions).

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- It is shown in Chen and Xia (2023) that the theory of univariate mixed trigonometric-polynomials (MTP) is **surprisingly decidable**.

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- Hence, a **reduction** from the unbounded case to the bounded case should be **favorable**.

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- Theorem (ISSAC '24, Chen & Xia Thm. 3.1)

Let $f_i \in S[\sin x, \cos x]$ for $i = 1, \dots, s$, then there are effective bounds $N, M \in \mathbb{R}$ such that any quantifier-free formula $\varphi(x)$ whose atoms are of the form $f_i \triangleright 0$ is true for all $x \in \mathbb{R}$ if and only if $\varphi(x)$ is true for all $x \in [N, M]$.

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- Our ISSAC'23 result can be directly derived from this by setting $S = \mathbb{Q}[x]$.

Applications of the Reduction Theorem

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As a corollary, the **reachability problem** of a linear differential system $\dot{x}(t) = Ax(t)$ is decidable, if the imaginary part of **eigenvalues** of A spans a **1-dimensional space** over \mathbb{Q} .

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Unifying $\sin x$ and $\cos x$ by substitution

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$$\sin x \mapsto \frac{2 \tan \frac{x}{2}}{1+\tan^2 \frac{x}{2}}, \quad \cos x \mapsto \frac{1-\tan^2 \frac{x}{2}}{1+\tan^2 \frac{x}{2}}.$$

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- Notice that

$$f \triangleright 0 \Leftrightarrow \sigma(f) \triangleright 0 \Leftrightarrow (1 + \tan^2 \frac{x}{2})^\ell \sigma(f) \triangleright 0.$$

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- The **singularity** of $\tan \frac{x}{2}$ at $x = (2k+1)\pi$ ($k \in \mathbb{Z}$) can be **treated separately**.

Essential Ingredient From Real Algebraic Geometry

Collins (1975) proposed Cylindrical Algebraic Decomposition (CAD). The key theorem in the paper is:

Theorem (Col75, Thm. 1)

Let $f(\vec{x}; y)$ be a **parametric univariate polynomial** and let C be a connected parameter region.

Suppose $\text{LC}_y(f) \neq 0$ for all $\vec{x} \in C$ and the **number of distinct complex roots** of $f(\vec{x}; y)$ is **invariant** for all $\vec{x} \in C$.

Then f is **delineable** over C , i.e. the real roots of f are **continuous functions** in the parameters.

He counted the complex roots by **his subresultant theory**.

Essential Ingredient From Real Algebraic Geometry (cont'd)

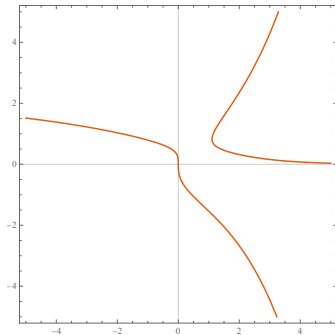


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 $f(x, y) = y^3 - (e^x - 1)y + x$

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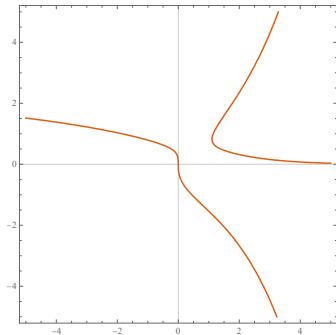


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① $\Delta_0 =$
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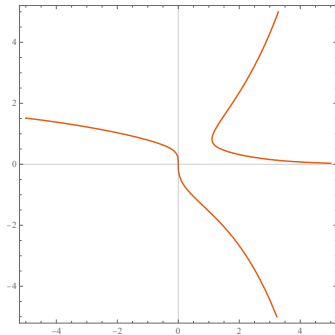


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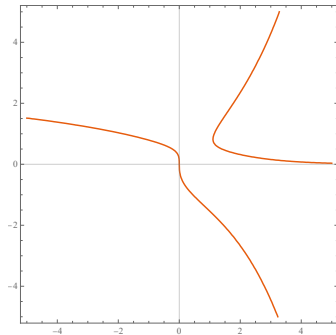


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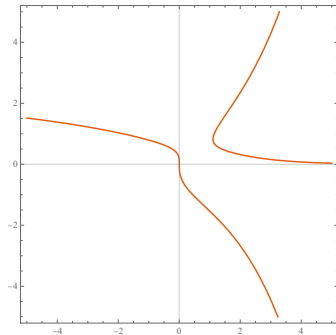


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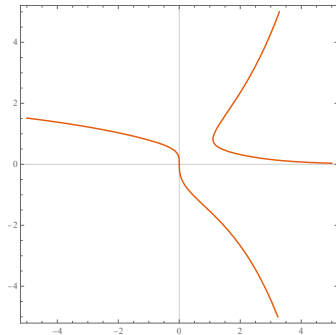


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Essential Ingredient From Real Algebraic Geometry (cont'd)

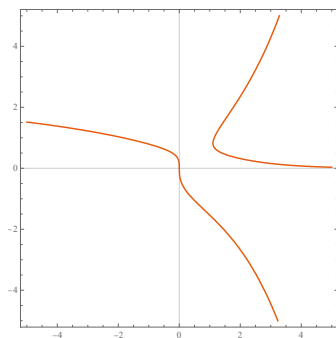


Figure 2: The locus of
 $f(x, y) = y^3 - (e^x - 1)y + x$

- The sub-discriminants:
 - ① $\Delta_0 = 27x^2 - 12e^x + 12e^{2x} - 4e^{3x} + 4$;
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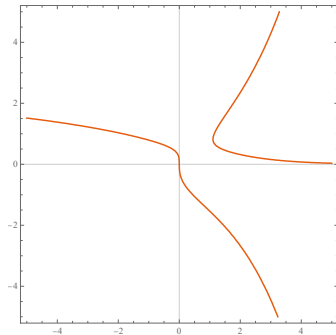


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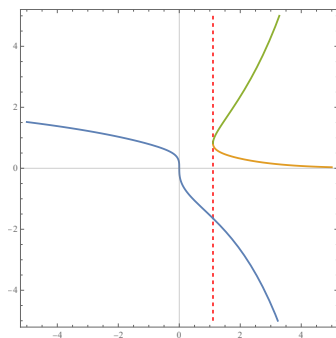


Figure 2: $f(x, y)$ is delineable when $x > 1.105\dots$

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- For all $x_0 > 1.105\dots$, $\Delta_0 \neq 0$, thus the real roots are continuous functions $y_1(x)$, $y_2(x)$ and $y_3(x)$.

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- e.g. $\tilde{\varphi}(x, y) = \text{“}y^3 - (e^x - 1)y + x < 0\text{”}$.
- Note that the subdiscriminants $\Delta_0, \Delta_1, \dots$ of $h \in S[y]$ is always in S , so they have finitely many real roots if they are not identically zero.

The Proof for the Reduction Theorem (cont'd)

- Then there is an interval $I = (M, +\infty)$ such that h is delineable over I .

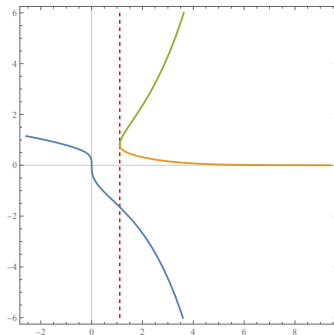


Figure 3: $h(x, y)$ is delineable over $I = (1.105\dots, +\infty)$

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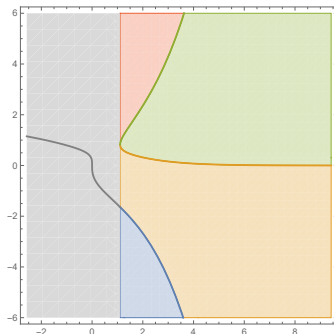


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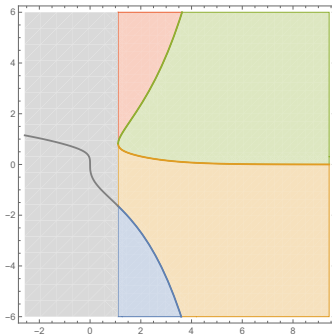


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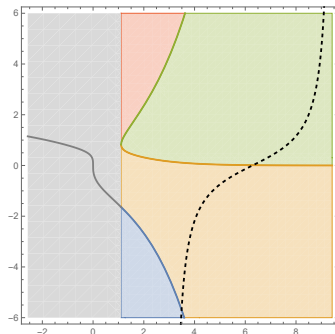


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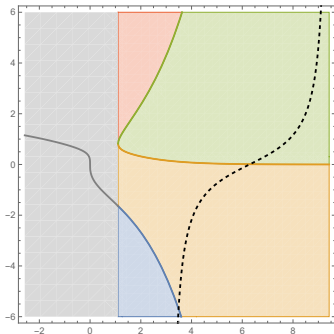


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- It is easy to observe that the graph of $y = \tan \frac{x}{2}$ intersects with each cell in each period.
- Hence, if $\tilde{\varphi}(x, y)$ holds in cell C , $\varphi = \tilde{\varphi}(x, \tan \frac{x}{2})$ is also satisfiable in each period $(2k\pi - \pi, 2k\pi + \pi) \subseteq I$. It suffices to look at one of them.

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4 Proof Sketch

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- Hence, our result in the univariate case is **not very far from being optimal**.

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Package TranscendentalProblems

We implement the reduction algorithm with Mathematica 13. Our package TranscendentalProblems is available at:

<https://github.com/xiaxueqq/TranscendentalProblems>.

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mission passed!

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respect +

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Experiments

We report some experimental data here.

Examples	1	2	3	4	5	6
Time (s)	1.531	0.969	1.797	3.562	1.891	1.438
Examples	7	8	9	10	11	12
Time (s)	0.125	0.031	0.188	0.078	8.406	2.109

Table 1: Running Time

Thank you!

You are more than welcome to give any suggestion!